# On the existence of a continuum of logics in $NEXT(\mathbf{KTB} \oplus \Box^2 p \to \Box^3 p)$

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#### Abstract

In this paper we consider formulas in one variable in the normal logic  $\mathbf{T}_2 = \mathbf{KTB} \oplus \Box^2 p \to \Box^3 p$ . Next, we use the formulas to define a continuum of logics over  $\mathbf{T}_2$ .

# 1 Introduction

In this paper we investigate normal modal logics over  $\mathbf{T}_2 = \mathbf{KTB} \oplus \Box^2 p \to \Box^3 p$ . The **KTB** logic is known as the Brouwer system and is an example of a non-transitive logic. It is characterized by the class of reflexive symmetric and non-transitive frames. There is very few results concerning this logic; some of them are included in [1] - [5].

Let us notice that adding the axiom  $\Box^2 p \to \Box^3 p$  to the Brouwer logic involves the following first order condition on frames:

$$(tran_2) \quad \forall_{x,y} (\text{if } xR^3y \text{ then } xR^2y).$$

The above property is known as a two-step transitivity.

# 2 Formulas in one variable in $NEXT(T_2)$

In this section we remind ourselves the infinite sequence of non-equivalent formulas in one variable defined in [2]. Denote  $\alpha := p \land \neg \Diamond \Box p$ .

Definition 1.

$$\begin{split} A_1 &:= \neg p \land \Box \neg \alpha, \\ A_2 &:= \neg p \land \neg A_1 \land \Diamond A_1, \\ A_3 &:= \alpha \land \Diamond A_2, \end{split}$$

For  $n \geq 2$ :

$$A_{2n} := \neg p \land \Diamond A_{2n-1} \land \neg A_{2n-2},$$
  
$$A_{2n+1} := \alpha \land \Diamond A_{2n} \land \neg A_{2n-1}.$$

Let us define the following model (see Figure 1):

**Definition 2.**  $\mathfrak{M} = \langle W, R, V \rangle$ , where

$$W := \{x_1, x_2\} \cup \{y_i, i \ge 1\}.$$

The relation R is reflexive, symmetric, 2-step accessible and:

$$\begin{array}{ll} x_1Rx_2, & x_2Ry_i \ \ for \ any \ \ i \geq 1, \\ y_iRy_j \ \ if \ \ |i-j| \leq 1 \ \ for \ any \ \ i,j \geq 1. \end{array}$$

The valuation is the following:

$$V(p) := \{x_1, x_2\} \cup \{y_{2m+1}, m \ge 1\}.$$

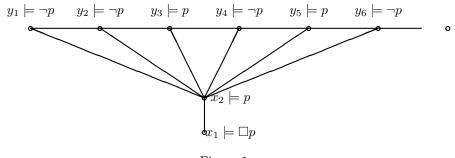


Figure 1.

**Lemma 3.** For any  $i \ge 1$  and for any  $x \in W$  the following holds:

$$x \models A_i \quad iff \quad x = y_i.$$

*Proof.* A detailed proof is presented in [2].

**Theorem 4.** The formulas  $\{A_i\}, i \geq 1$  are non-equivalent in the logic  $\mathbf{T}_2$ .

Proof. Obvious.

# 3 The existence of a continuum of logics over $T_2$

Yutaka Miyazaki in [4] proved the existence of a continuum of logics over **KTB**. First, he showed the existence of a continuum of orthologics and then applied an embedding from orthologics to **KTB** logics. In [5] Y. Miyazaki proved the existence of a continuum of logics over  $T_2$ . In the proof he considered logics determined by the so-called wheel frames.

**Definition 5.** For  $n \in \omega$ ,  $n \geq 5$ , the wheel frame  $\mathfrak{W}_n = \langle W, R \rangle$  of degree n consists of the following set and binary relation:  $W = rim(W) \cup h$ , where rim(W) := $\{1, 2, ..., n\}$  and  $h \notin rim(W)$ . Any element in rim(W) is called a rim element, whereas the element h - the hub element. The relation R is defined as: R := $\{(x, y) \in (rim(W))^2 : |x - y| \leq 1(mod(n - 1))\} \cup \{(h, h)\} \cup \{(h, x), (x, h) : x \in$  $rim(W)\}.$ 

For example in Figure 2 we present a diagram of the  $\mathfrak{W}_8$ .

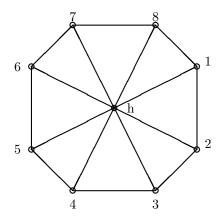


Figure 2.

Y.Miyazaki proved the following two lemmas (Proposition 19, Lemma 20 from [5]):

**Lemma 6.** For  $m > n \ge 5$ ,  $L(\mathfrak{W}_n) \not\subseteq L(\mathfrak{W}_m)$ .

**Lemma 7.** For  $m \ge n \ge 5$ , suppose there is a p-morphism from  $\mathfrak{W}_m$  to  $\mathfrak{W}_n$ . Then m is divisible by n.

On the base of these two lemmas and by using the splitting technique effectively, Y. Miyazaki constructed a continuum of normal modal logics over  $\mathbf{KTB} \oplus \Box^2 p \to \Box^3 p$  logic.

Below, we present an improved method (in comparison to the one from [2]) for obtaining a continuum of logics above  $\mathbf{T}_2$ . Actually, we axiomatize the logics determined by wheel frames with formulas in one variable. Let us define new formulas for  $k \in \omega$ :

### Definition 8.

$$C_k := \Box^2[A_{k-1} \to \Diamond A_k], \text{ for } k > 2,$$
  

$$D_k := \Box^2[(A_k \land \neg \Diamond A_{k+1}) \to \Diamond \varepsilon],$$
  

$$E := \Box^2(\Box p \to \Diamond \gamma),$$
  

$$F_k := (\Box p \land \bigwedge_{i=2}^{k-1} C_i \land D_{k-1} \land E) \to \Diamond^2 A_k,$$

where

**Lemma 9.** Let  $k \ge 5$  and k be an odd number. Then  $\mathfrak{W}_i \not\models F_k$  iff i is divisible by k+2.

Proof. ( $\Leftarrow$ ) Let i = k + 2. We define the following valuation in the frame  $\mathfrak{W}_i$ :

$$\begin{array}{rrrrr} h & \models & p, \\ 1 & \models & p, \\ 2 & \models & p, \\ 3 & \not\models & p, \\ 4 & \not\models & p, \\ 2n-1 & \models & p, \text{ for } n \geq 3 \text{ and } 2n-1 \leq i, \\ 2n & \not\models & p, \text{ for } n \geq 3 \text{ and } 2n < i. \end{array}$$

Then the point 1 is the only point such that  $1 \models \Box p$ . And further:

$$\begin{array}{rcl}
h & \models & p, \\
2 & \models & \gamma, \\
3 & \models & A_1, \\
4 & \models & A_2, \\
k+1 & \models & A_{k-1}, \\
k+2 & \not\models & A_k, \text{ and } k+2 \models \varepsilon
\end{array}$$

Then we see that for all j = 3, ..., k + 1 we have:  $j \models A_n$  iff n = j - 2. We conclude that for all j = 1, 2, ..., k + 1 it holds that:  $j \models \bigwedge_{i=2}^{k-1} C_i \wedge D_{k-1} \wedge E$ . Then the predecessor of the formula  $F_k$ :  $(\Box p \land \bigwedge_{i=2}^{k-1} C_i \land D_{k-1} \land E)$  is true only at the point 1. At the point 1 we also have:  $1 \not\models \diamondsuit^2 A_k$ , because there is no point in the frame satisfying  $A_k$ . Hence at the point 1, the formula  $F_k$  is not true.

In the case when i = m(k+2) for some  $m \neq 1, m \in \omega$  we define the valuation similarly:

for all l such that:  $0 \le l < m$ . The rest of the proof in this case proceeds analogously to the case i = k + 2.

 $(\Rightarrow)$  Suppose there is a point  $x \in W$  such that:

$$\begin{array}{lll}
x & \models & (\Box p \wedge \bigwedge_{i=2}^{k-1} C_i \wedge D_{k-1} \wedge E) \\
x & \models & \neg \diamondsuit^2 A_k.
\end{array}$$

First, let us observe that  $x \neq h$  because if  $h \models \Box p$  then for all  $x \in rim(W)$  we have  $x \models p$ . Hence there is no point  $x' \in rim(W)$  such that  $x' \models \gamma$ . Then it is impossible that at the point h formula  $(\Box p \land \bigwedge_{i=2}^{k-1} C_i \land D_{k-1} \land E)$  is true. Then x has to belong to rim(W). Let x = 1. Then we know that there is a point 2 such that  $2 \models \gamma$  what involves existence of the next point 3 such that  $3 \models A_1$ . Because of  $C_i$ , i = 1, 2, ..., k - 1 we know that there is a sequence of points 3, 4, ..., k + 1

such that  $n \models A_{n-2}$  for  $2 \le n \le k+1$  and  $k+1 \models \neg \Diamond A_k$ . Then a point k+2, which is next to the point k+1, has to validate the formula  $\varepsilon$ . Because  $h \not\models \varepsilon$  and  $k, k+1 \not\models \varepsilon$  then it must be a new rim element. It has to see some point validating  $\Box p$  and if it sees the point 1 then we have that i = k+2. But suppose that k+2does not see the point 1. Anyway, it has to see another point validating  $\Box p$ . Say, it is the point k+3. But it has to be  $k+3 \models \Diamond \gamma$ . Because  $h \not\models \gamma$  then it has to be other point, say k+4 such that  $k+4 \models \gamma$ . Then there has to be a next point k+5 different from h such that  $k+5 \models A_1$ . Again from  $C_i$  for i = 1, 2, ..., k-1we have to have:  $k+6 \models A_2, ..., 2k+3 \models A_{k-1}$ . Then we have that there has to be a point 2k+4 validating  $\varepsilon$ , and then some point validating  $\Box p$ . If it is the point 1 then we have i = 2(k+2). If not, then we have analogously another sequence of k+2 points and so on.

The main theorem is the following:

**Theorem 10.** There is a continuum of normal modal logics over  $\mathbf{T}_2$  defined by formulas written in one variable.

*Proof.* Let  $Prim := \{n \in \omega : n+2 \text{ is prime}, n \geq 5\}$ . Let  $X, Y \subset Prim$  and  $X \neq Y$ . (Exactly:  $X \not\subseteq Y$  and  $Y \not\subseteq X$ ). Consider logics:  $L_X := \mathbf{T_2} \oplus \{F_k : k \in X\}$  and  $L_Y := \mathbf{T_2} \oplus \{F_k : k \in Y\}$ . Let  $j \in Y \setminus X$ . Obviously:  $F_j \in L_Y$ . From Lemma 9 we know that  $F_j \notin L_X$ , because  $W_{j+2} \not\models F_j$  and  $\forall_{i \in X} [i \neq j \Rightarrow W_{j+2} \models F_i]$ . That means that we are able to define a continuum of different logics above  $\mathbf{T_2}$  by formulas of one variable.

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